



Computing evolution of ATLAS and CMS

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Outline

- **ATLAS** computing upgrades
- **-CMS** computing upgrades



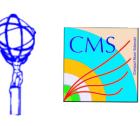




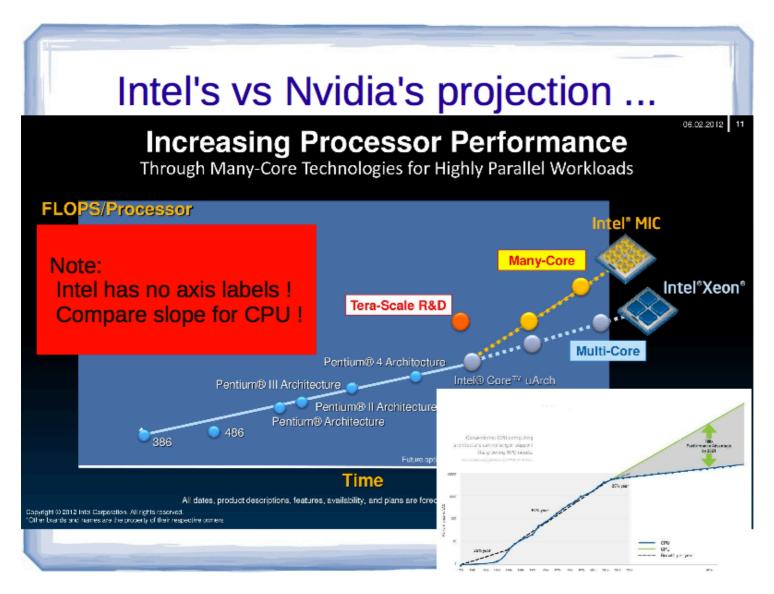
Why we need upgrades to our computing models?

- · We can identify at least 5 drivers for upgrade computing
 - 1. consequences of increasing pile-up, event complexity and size
 - 2. consequences of new detectors and triggers
 - 3. consequences of increasing sample size
 - 4. consequences of new architectures
 - 5. consequences of new software technologies





New architectures of Processors



- Going towards many-core architectures
 - Emerging consensus that to profit in future from performance increase code must utilize many core platforms
 - ATLAS and CMS software will have to be modified to fit new CPU platforms
- Technology may develop faster than we expect
 - Industry may require us for best performance to go to many/multi core already before Phase-1 (2014...)







Implications of architectures' evolution

- The frameworks are already adapting to new architectures
- Different approaches
 - WholeNode
 - preferred by CMS (job handles the machine)
 - Multicore
 - generally preferred by site admins to keep the sites busy
 - we will probably converge on this one
- Beyond this, the software algorithms must also adapt (reconstruction, simulation)
 - · Concurrency, parallel programming
- The distributed computing must also adapt to make efficient use (e.g. whole-node scheduling)

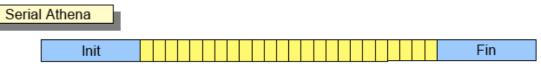


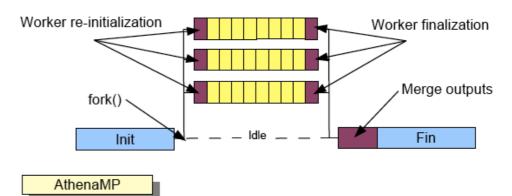




Multiprocessing in ATLAS

- Current approach
 - ·A variation of the single-thread framework (Athena), called AthenaMP, is being rolled-out
 - · Athena uses N threads
 - · Events are split into blocks
 - · Each block is processed by a separate AthenaMP thread (workers)
 - · A common thread collects the results and handles the output (finalize)
 - ·Helps control the memory issues, but has limitations by ~32 cores
 - Synergies: AthenaMP being investigated for HLT
 - Extensive work on IO new framework planned
 - A re-write of the Gaudi core is also likely
 - Process management migrated h Python multiprocessing to a custom C++ library





- Actively working to adapt the software algorithms too
 - Efficiency and architectural work is starting
 - Optimization of existing code
 - Exploring GPGPUs, especially for tracking

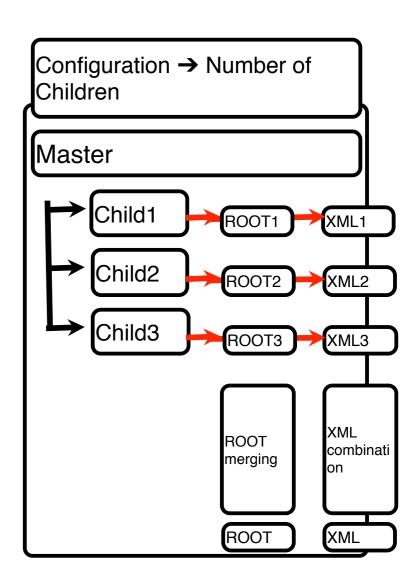






Multiprocessing in CMS

- Current implementation uses the cores via fork
 - Relying on the Kernel COW to allocate only once the shared memory payloads
- JobWrapper configures the number of children
 - Either via workflow settings (Manycore)
 - Or using /proc/cpuinfo to use the whole node
- JobWrapper executes a single CMSSW job producing master xml file and multiple FrameworkJobReport.xml and output files
 - · Like MyAnalysis_0.root, MyAnalysis_1.root, ... MyAnalysis_N.root
- JobWrapper merges all ROOT files and stages it out to MSS and also combines all xml into one







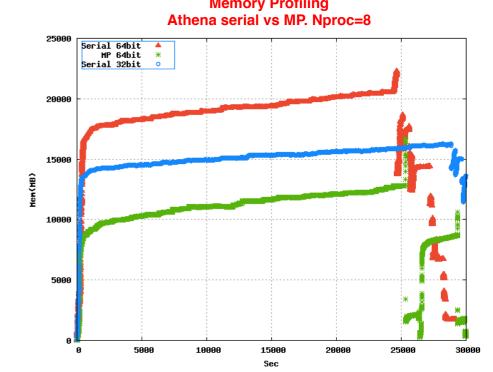
Multiprocessing performance in ATLAS

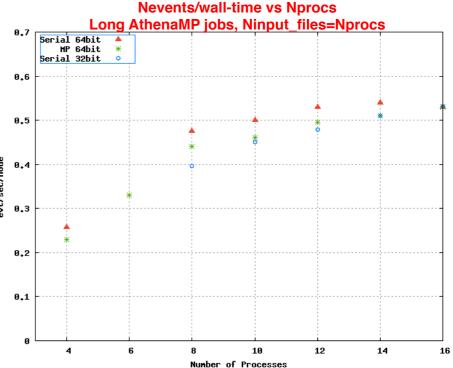
- Saves up to 40% memory per core (RSS)
 - Scales well with cores, tested up to 32 forked processes
- A single AthenaMP job with N workers of course runs faster than the corresponding serial job, but not N times faster
 - Amdahl's law:

$$S(N) = \frac{1}{(1-P) + \frac{P}{N}}$$

 \boldsymbol{P} is the portion of the program that can be made parallel $\boldsymbol{S(N)}$ the maximum speedup that can be achieved by using \boldsymbol{N} processors

- Long reconstruction jobs result in large output files
 - With high number of workers it can become problematic to manage such big files
 - Also, the output validation of such files is slow
 - In order to address this issue we would need to have a mechanism, which allows single job make more than one output file of a given type





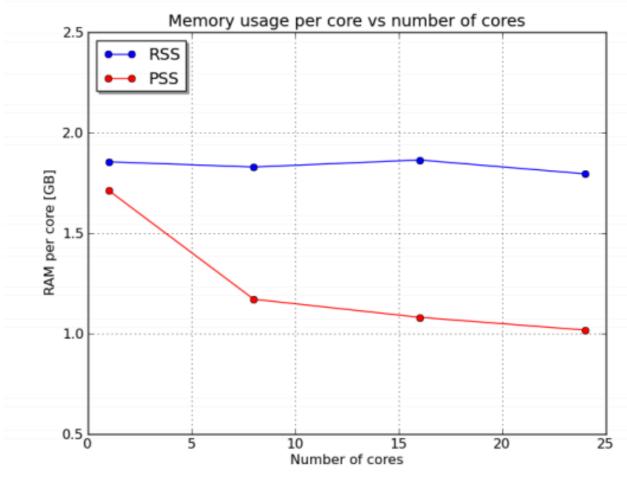






Multiprocessing performance in CMS

- · Using a reasonable metric (VSIZE is not, and not even RSS), up to 40% memory gained per core
 - Scales well with cores, tested up to 96 forked processes
- Price to pay is a <5% CPU inefficiency
 - Not due to processing itself, which keeps all the CPUs at 100%, but
 - Skew between forked processes: they have to wait for the last before finishing
 - Time spent in merging results is idle for all the cores
 - Becomes ~irrelevant for jobs lasting more than 4 hours



· As of today:

- Multicore approach simply works (it is in operation on specific T1s queues)
- Not too much pressure to use it as standard, since CMS fits quite well with the 2 GB/core limit (after having spent 2011 to reduce the RAM footprint even in presence of PileUp ~ 40)







Is forking enough?

- Up to a certain point
 - It helps to reduce memory footprint by sharing common payloads
 - But it relies on processing N events in parallel
 - •When events are big, the "not shareable event data" can choke systems

After that

- We need parallelism within the single event
- Smaller memory footprint: a single event in memory
- But: parallelism at event level→ we need to touch the algorithms





Work in progress for the software [1]

- Short, medium and long-term work necessary
 - Common issues to be addressed
 - Memory saving to be able to keep all cores in a machine busy (the usual, in principle trivial parallelism on event level)
 - Better usage of the resources

· ATLAS

- · I/O framework, getting rid of POOL, etc advancing well
- Full usage of each core: vectorizing a few algorithms
 - Tracking will be explored first
 - Likely influence on data model
- Parallelism on intermediate levels: on algorithm level, on sub-event data level. Both require refined communication mechanisms
- Common work with PH-SFT strongly considered. Common tracking effort with CMS desirable







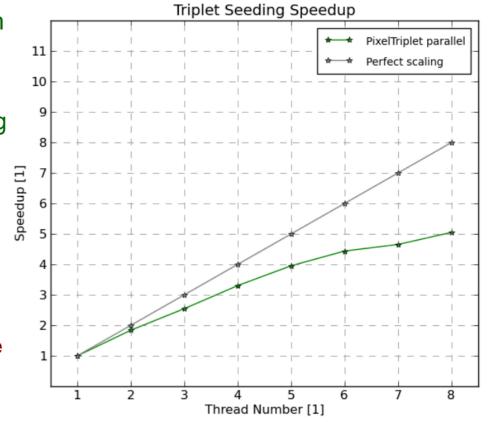
Work in progress for the software [2]

- CMS

- · libdispatch (Apple): send atomic tasks to a global queue, which dispatches it to the cores when available
 - Close to a batch system handling thread dispatching in the system
 - Not (too) exposed to physicist: can be hidden at FW level, we do not need to be all multi-threading experts
 - CMSSW cores become tasks in the queue, FW to resolve ordering

Use Intel Threading Building Blocks (TBB)

- Test application on a single algorithm:
 CMS tracker pixel seeding
 - Which is just a loop over hit triplets to see whether they fit an helix
 - Seed candidates are divided into blocks, and sent to different threads
 - With 8 threads, seeding goes from 14% of total reconstruction time to virtually negligible
 - RAM penalty is very small (~ 1 MB/thread)
 - Much better than running pixel seeding on 8 different events as in forking







Work in progress for the software [3]

Vectorization

- Up to here we were just trying to use in a better way multi core CPUs
- We can squeeze more performance by using the vector units which are paired to them (MMX, SSE, SSE2, ..., AVX, ...)
- ATLAS and CMS are experimenting vectorization techniques for their software
- Step #1 (CMS): use auto-vectorization in latest GCC; studies ongoing



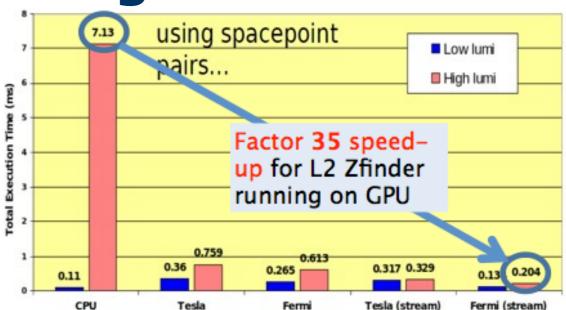


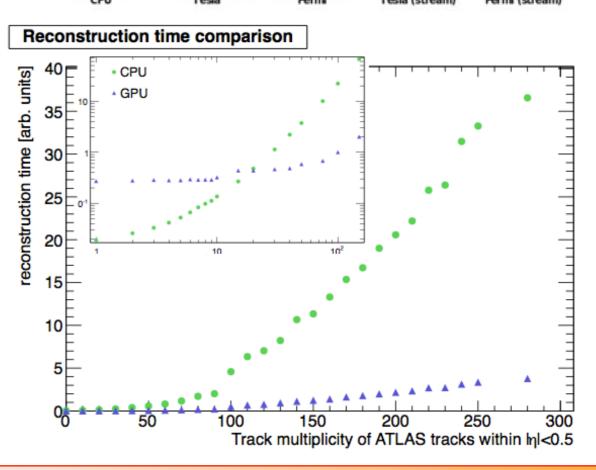


ATLAS GPU-based Tracking

- First tracking prototypes for Level-2 track trigger and offline tracking
 - concentrate on aspects of track reconstruction chain
 - z-vertex finder
 - track seed finder
 - Kalman filter
 - · early phase, still significant approximations
- very significant timing gains
 - but: lots of software development needed to obtain precise tracking
 - · investigate mixed scenario?
 - e.g. combinatorial seed finder on GPUs
 - CPUs for serial processing steps to do precision calculations

~150 speed-up seen in other Kalman filter studies Experience with GPUs can help with many-cores











Upgrading the ATLAS Computing Model

- The ATLAS Distributed Computing & the Grid are doing very well.
 - 1000's of users can process petabytes of data with millions or more of jobs
- · But at the same time, we are starting to hit some limits:
 - Scaling up, elastic resource usage, global access to data
- What can we learn from external innovations? (without disrupting operations!)
- Various R&D Projects and Task Forces were formed one year ago
 - NoSQL databases R&D
 - Cloud Computing R&D
 - **XROOTD Federation and File level Caching Task force**
 - **Event Level Caching R&D**
 - Tier3 Monitoring Task Force
 - **CVMFS Task Force**
 - **Multicores Task Force**
 - ·Also Network Monitoring...
- Deploying and using LHCONE
 - ·Already done in some sites, no immediate gain but almost transparent migration
- https://twiki.cern.ch/twiki/bin/viewauth/Atlas/TaskForcesAndRnD







Upgrading the CMS Computing Model

- Job Submission
 - Going mostly to Glideln WMS
- Storage
 - Differentiate T1D0 and T0D1 to help analysis @ Tier1s
 - No wild tape recall
- Remote (WAN) data access
 - US and EU level projects are effectively building up "federation" of centers
- Software
 - Redesign of critical parts to exploit multi core systems more efficiently; going beyond with parallelization
- Networking
 - LHCONE to be deployed soon (indeed, already partially done)







Short and long term work in LHC

- Short and medium-term work well under way
 - Important technologies: Cloud computing, Hadoop basket
 - · Common work ATLAS/CMS/IT-ES on job management with Panda+Condor
- Long term
 - · Future role of middleware?
 - Try consolidate middleware flavors
 - Possible consequences for systems we have on-top
 - · Or rather try to be independent of middleware
- Common solutions
 - · With other experiments, with CERN (IT, PH), with other labs
 - · Many of the TEG areas are chances for commonality
 - Concrete progress in a few areas so far
 - ATLAS, CMS, IT-ES: common analysis framework based on PanDA+Condor/Glidein
 - Helix/Nebula: Cloud computing project involving CERN/ATLAS, EMBL, ESA, and 13 industrial partners
 - Other opportunities
 - Storage federations, network monitoring, data preservation
- Common solutions are needed when manpower/funding shrink (EGI, EMI deadlines, OSG cuts)







Common Analysis framework ATLAS/CMS

- Initiative from CERN IT-ES, ATLAS and CMS for a common analysis framework started March 13th 2012
- -Assess the potential for using common components for distributed analysis, based on elements from PanDA and glidelnWMS
- Initial plan
 - Feasibility study Mandate: http://cern.ch/go/9mNC
 - Analyze architectures of both experiment's analysis frameworks
 - Identify interfaces to external systems
 - Identify what can be reused
 - How much effort is it?
 - Identify show-stoppers
 - Functionality study
 - What do ATLAS and CMS gain and loose in terms of functionality by adopting a common framework
 - Operations study
 - What is the impact on the cost of operating various proposals
- A common analysis framework could lead the way to further commonalities and collaboration between the experiments in the future







Virtualization and Cloud R&D in ATLAS

- Active participation, almost 10 persons working part time on various topics
 - · https://twiki.cern.ch/twiki/bin/view/Atlas/CloudcomputingRnD
- Data Processing
 - Panda Queues in the Cloud
 - Centrally managed
 - Tier3 Analysis Clusters (Instant Cloud Site)
 - User/Institute Managed, Low/Medium Complexity
 - Personal Analysis Queue (~One click, run my jobs)
 - User Managed, Low Complexity (almost transparent)
- Data Storage
 - · Short term data caching to accelerate above data processing use cases
 - Transient data
 - Long term data storage in the cloud
 - Integrate with DDM



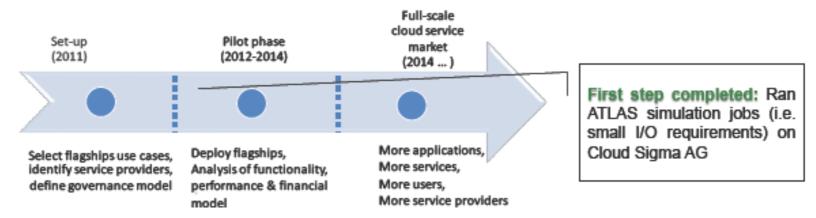




Cloud Computing in ATLAS

- Helix Nebula: the Science Cloud
 - European Cloud Computing Initiative: CERN, EMBL, ESA + European IT industry
 - Evaluate cloud computing for science and build a sustainable European cloud computing infrastructure
 - Set up a cloud computing infrastructure for the European Research Area
 - Identify and adopt policies for trust, security and privacy at a European-level
 - CERN/ATLAS is one of three flagship users to test a few commercial cloud providers (Cloud Sigma, T-Systems, ATOS...)





- Simple approach
 - **Use CernVM with preinstalled SW**
 - Condor pool with master at CERN (one of the pilot factories)
 - I/O copied over the WAN from CERN (lcg-cp/ lcg-cr for input/output)

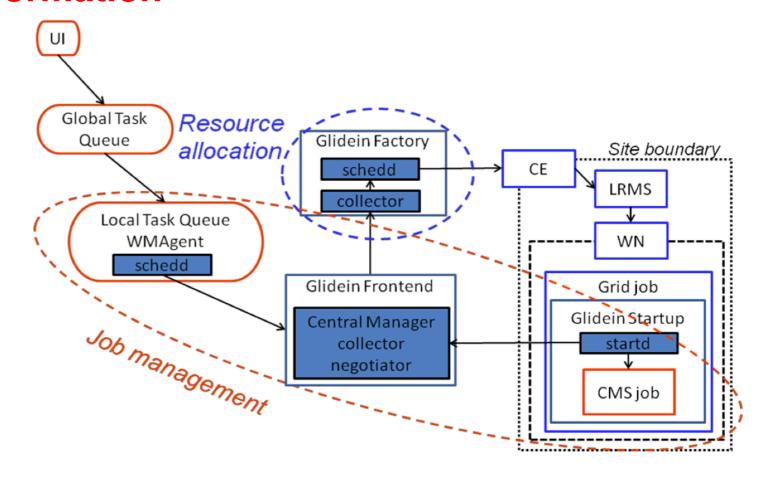






Job Management in CMS

- ·Use of the glidein WMS factory to access WN slots.
- No reliance on BDII load information
- Has proven to be able to sustain > 50k jobs in production
- Having a single Global queue allows for CMS wide prioritization









Storage issues in CMS

- The main issue for hosting physics analysis @ Tier1s is the risk of having chaotic file recalls from Tape
- · While CMS up to now has envisaged a flat T1D0 disk model, with all the files logically "living on tape", a split with some files pinned to disk (T0D1) is the preferred solution for analysis use cases
 - Being investigated
- Streaming access to data
 - One of the consequences of a next generation networking (GARR-X) is the possibility to at least partially overcome the "jobs go to where data is" paradigm
 - With current CMSSW on AOD events, it is typical to see analysis jobs reading < 500 KB/s from storage. With 10 Gbps, hundreds of job reading over WAN are possible (using vector reads)







Which storage model?

- Common problem in ATLAS and CMS
- Driving ideas
 - We cannot avoid completely data placement too early
 - The current paradigm "jobs go where data is" stays valid in the vast majority of the cases
 - But a number of use cases can be better served by streaming:
 - 1.Fallback in case of hardware problems: In case of unavailability of the local copy of a file the user job can transparently access a remote replica, thus increasing the processing efficiency and decreasing job failure and resubmission
 - 2. Specific interactive use cases with low I/O processes, i.e. visualization programs.
 - 3. Address site congestion, when the available copies of a dataset are at overloaded sites (overflow)
 - 4.Increase the utilization of CPU power at sites where proper data management is not possible (for example, small University sites with no/small storage)

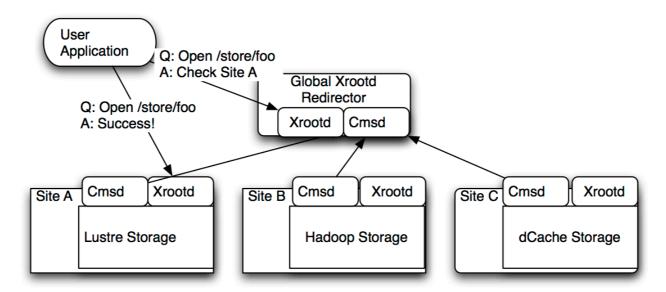






Streaming data: Storage federations

- Why storage federations?
 - · Fitting the experiments' needs
 - Little needed for code development
 - Excellent access to Xrootd team if development is needed
 - Mature product used for many years
 - Common technology with LHC experiments
 - ATLAS, CMS and Alice
 - Very efficient at file discovery
 - Works seeamlessly with the root data formats that we use
 - ROOT team collaborating with Xrootd collaboration on efficient wide area data access







Example: storage federations status in ATLAS

- In US cloud, Tier 1 site and 80% (4 out of 5) Tier 2 sites part of Federation
 - Soon all Tier 2 sites will be part
- Using X509 authentication for reading
 - · Plugin code looks for ATLAS VOMS extension to allow access
 - Deployed at MWT2, AGLT2, SWT2 and SLAC T2 currently
 - The rest (BNL, and NET2) shortly
- Prototype Midwest region redirector setup at Chicago
 - · Prun jobs used to test regional redirector
- Work still needed
 - Integrate federated storage with job management system to combine the power of both
 - Change the Pilots to handle missing files at run time
 - Evolve Panda to understand federated storage and federated site queues

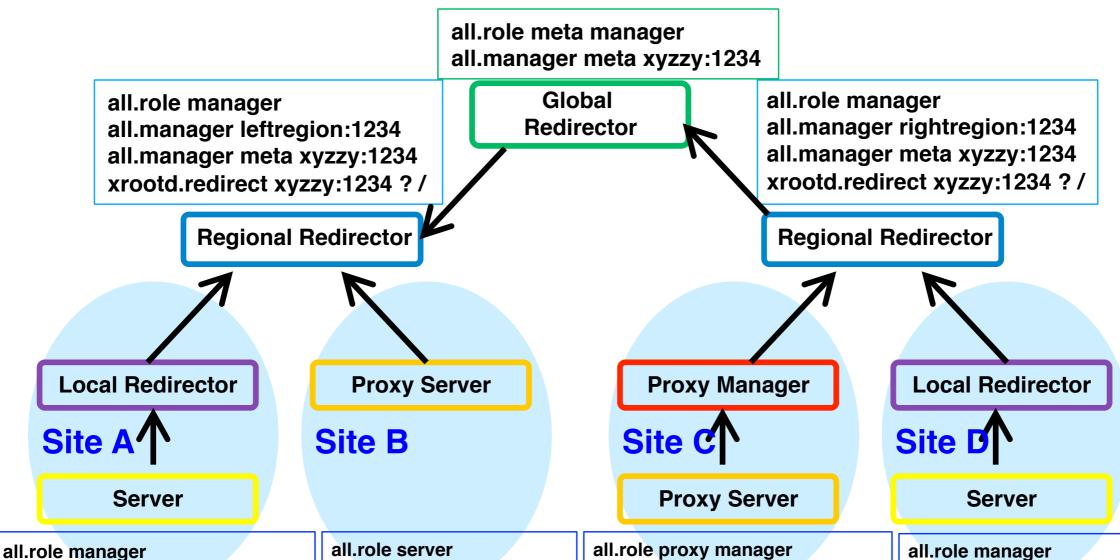
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Configured regional redirectors in ATLAS



all.manager lefthand:1234
all.manager meta leftregion:1234
xrootd.redirect leftregion:1234?/

all.role server all.manager leftregion: 1234 ofs.osslib psslib.so pss.origin mycluster:1094 all.role proxy manager all.manager meta leftregion: 1234 xrootd.redirect xyzzy:1234?/

all.role manager all.manager righthand:1234 all.manager meta rightregion:1234 xrootd.redirect rightregion:1234?/

Andrew Hanushevsky (SLAC)





Xrootd and CMS Software

Layered data access configured in **CMSSW**

- Regional redirectors already active '
- · UNL (Nebraska) for US sites
- Bari for EU sites
- · CERN will soon publish EOS/CAF via Xrootd

I need '/store/foo'

- I try local access (via 'direct' protocol) ... if not found
 - I try accessing a first level redirector (national, for example)... if not found

- I scale up to the global redirector... if not found
 - Sorry file not accessible

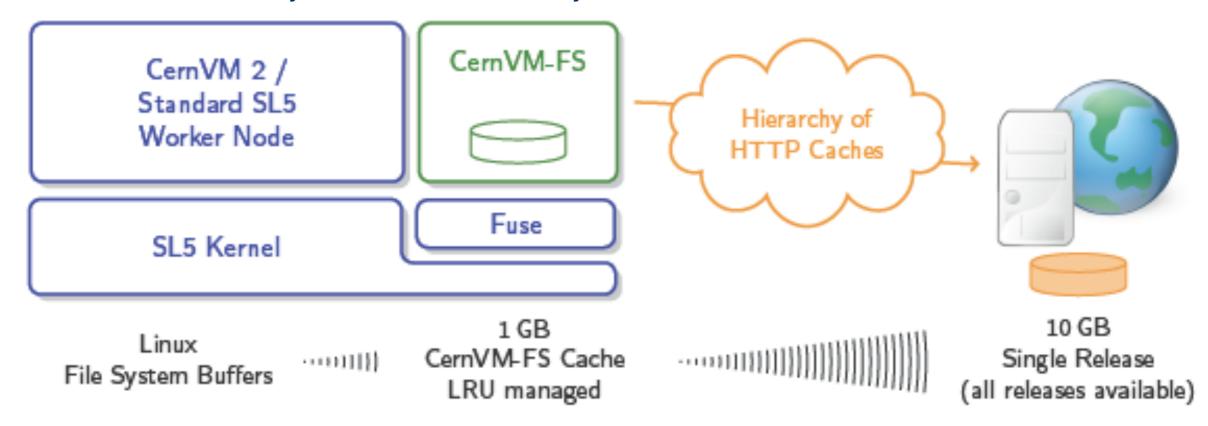






CVMFS

- · ATLAS is now fully using the dynamic software distribution model via CVMFS (CernVMFileSystem)
- ·Virtual software installation by means of an HTTP File System



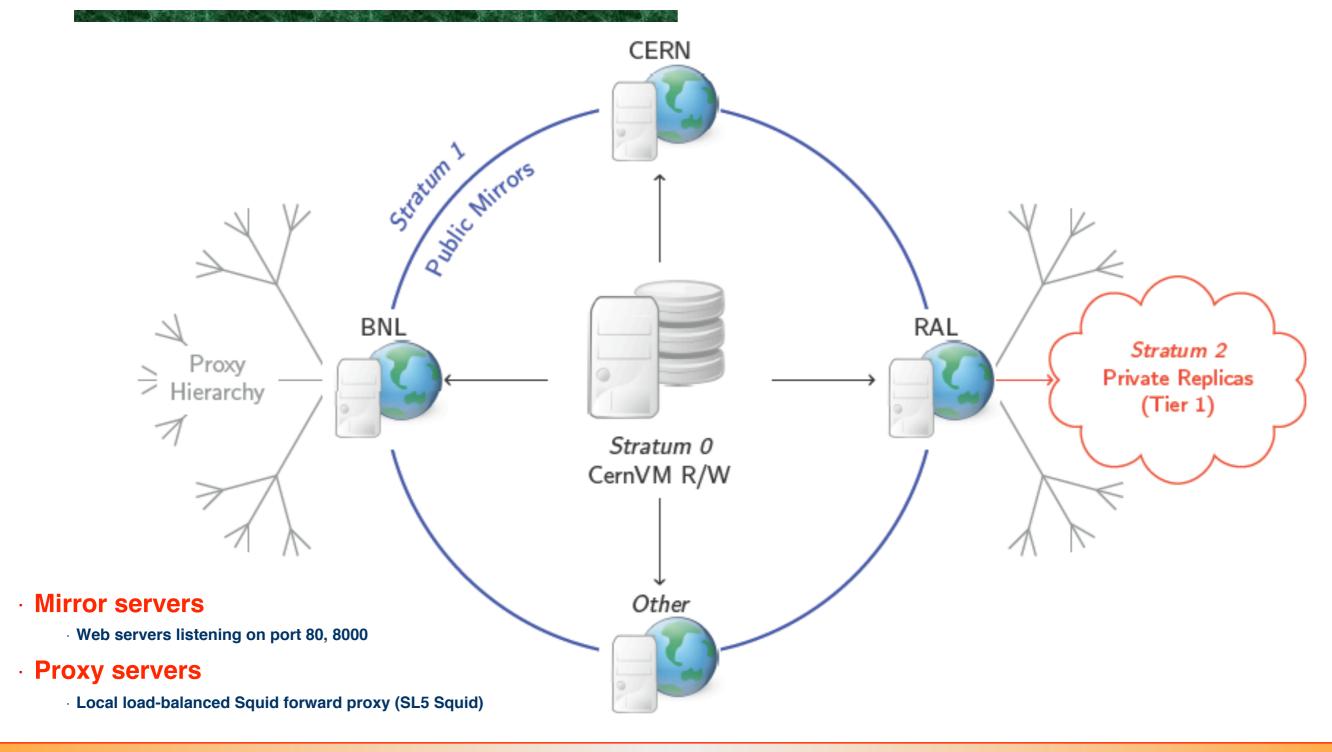
- Distribution of read-only binaries
- · Files and file meta data are downloaded on demand and locally cached
- · Self-contained (e.g./cvmfs/atlas.cern.ch), does not interfere with the base system
- CMS is moving from training phase to real use (# of sites increasing fast)







CVMFS Backends used by ATLAS









ATLAS and CVMFS

- Migration status of the ATLAS sites:
 - > 60% of the sites are now using CVMFS
 - The rest should migrate by the end of this year
- · /cvmfs/atlas.cern.ch
 - "Stable releases" software
 - New production server in CERN IT ready for production
 - Populated, but not yet in full production
- · /cvmfs/atlas-condb.cern.ch ATLAS Condition Flat Files
 - Release manager machine hosted by CERN IT
 - Automatic update several times a day
- · /cvmfs/atlas-nightlies.cern.ch ATLAS Nightlies
 - Tested on grid sites too
- Integrated with the current Installation System
 - · CVMFS sites are used by the installation system transparently, aside of sites using a different FS







Networking upgrades

- Tier2 sites are ready to switch to 10Gbps connectivity
 - · All of them have the appropriate hardware to handle that
 - We hope to have it soon, as in some cases we're suffering of link congestions (eg. Napoli), especially if the site already moved to LHCONE
- ATLAS in LHCONE
 - 3 out of 4 Tier2 sites are already connected
 - Napoli
 - Milano
 - Roma
- CMS in LHCONE
 - 3 CMS Tier2s out of 4 already had the transition to LHCONE
 - Bari: April 26th
 - Roma: May 7th
 - Pisa: May 10th
- No problems so far
 - The transition was done very quickly
 - No performance difference detected with respect to the previous setup







The ATLAS Computing model changes

- Move from a fully hierarchical model to less hierarchy and more Bandwidth
- 4 recurring themes:
 - Any site can replicate data from any other site
 - Multi Domain Production
 - Need to replicate output files to remote Tier-1
 - Dynamic data caching
 - Analysis sites receive datasets from any other site "on demand" based on usage pattern
 - Remote data access
 - Local jobs accessing data stored at remote sites
- ATLAS is now heavily relying on multi-domain networks and needs decent network monitoring
- Work Ongoing on global access/Data Federation



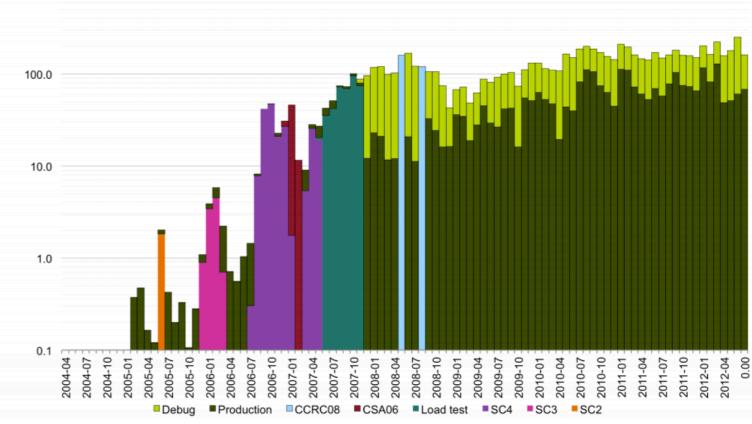


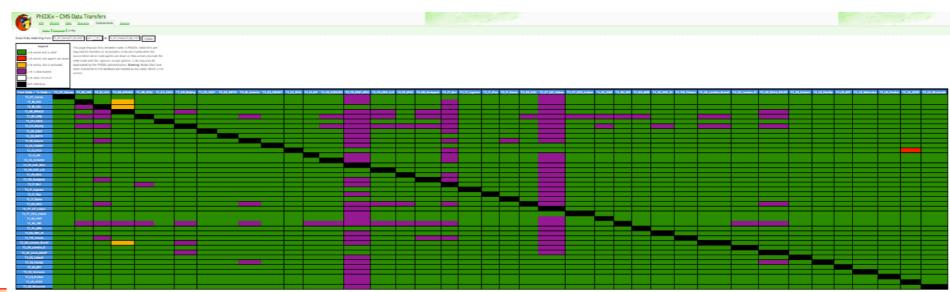


CMS - full mesh model

- Also CMS started data taking with a MONARC approach, fully hierarchical
- Already in 2010 this was changed to a model in which all the sites (most of ...) can speak together directly
- Situation end of 2010:

A part of the full matrix (T0,1,2) vs (T0,1,2): 2550 active links











Common Distributed Computing evolutions

- Database Scaling
 - Hadoop environment looks best
- Storage and data management
 - · Maintain stable storage for placed data
 - Support access from experiment jobs
- Workload management
 - · Pilots and frameworks
 - GlideinWMS
 - Whole node scheduling
- CPUs and I/O
 - Use of CPU affinity and pinning
 - · Handling of CPU-bound and I/O-bound jobs
- Exploring Common Solutions among several experiments/WLCG







Conclusions

- The ATLAS and CMS Computing models are about to evolve in the coming months in a variety of different aspects
 - Job handling and exploration of the Cloud Computing extensions
 - Storage models and data access
 - Data transfer model and data placement vs direct streaming
 - Usage of new CPU architectures and full exploitation of the multiprocessing
 - **Network upgrades**
- The Computing Infrastructure of ATLAS and CMS will continue working next year at full speed, taking advantage of the LHC stop to test and deploy new solutions
- "Common Solutions" seem more and more the way to proceed, while at the same time maintaining experiments' peculiarities

·Credits

- · Thanks to R. Jones, and D. Benjamin for the ATLAS material
- Thanks to C. Grandi and D. Bonacorsi for the CMS material